## Sequence Analysis for Life Course Data Modifying the OM Algorithm for better duration handling

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## Introduction

- Today I will present a new algorithm for comparing longitudinal trajectory data such as work or life histories
- A modification of the Optimal Matching algorithm, the dominant method in most current sequence analysis
- Very preliminary work, first test of the modified algorithm
- Thanks to comments from Cees Elzinga, and participants in the Geary Institute Seminar, April 2007.



## What is Sequence Analysis?

- Sequence Analysis (SA) treats sequences as units, and compares them holistically
- Sequences in this sense are longitudinal structures, e.g.,
  - dance (series of steps)
  - conversation (series of utterances)
  - macromolecules like DNA (series of CAGT "bases")
  - life course histories (series of time-units in life course state space)
  - that is, typically linear sequences of observations in a discrete state space
- Note the distinction between discrete sequences in discrete spaces and continuous-time sequences in discrete spaces: the latter we represent as sequences by discrete-ising time



## What is Sequence Analysis?

- SA works by defining pairwise distances between sequences according to some metric
- It typically proceeds by using the pairwise distances to generate data-driven typologies using cluster analysis
- Comparing sequences to reference sequences, and multidimensional scaling are also possible
- This holistic approach is an alternative to more conventional techniques which often focus on transition rates:
  - focus on the outcome (epiphenomenon?)
  - rather than the underlying generative processes



## Descriptive, exploratory

- Of descriptive, exploratory value rather than analytic/stochastic
- Can provide an digestible overview of complex longitudinal data
- Of particular promise with multi-domain data
- Or where generative processes are complex and changing (*e.g.*, along the lifecourse)
- However, when you want to test clearly specified hypotheses about the generative processes, conventional stochastic techniques such as hazard rate modelling are often much more powerful



## Defining pairwise distance

- Defining the pairwise distance is the foundation of SA
- Many possibilities:
  - Hamming:  $D_{AB} = \sum_{i Ai} B_i \text{ or } \sum_{i} d(A_i, B_i)$
  - Degenne:  $D_{AB} = \sum_{i} \cos^{-1}(\mathbf{X}_{Ai}, \mathbf{X}_{Bj})$  where  $\mathbf{X}_{Ai}$  the cumulated duration
  - Dijkstra–Taris (1995): delete repeats, delete non-common elements, count matches
  - Optimal matching algorithm (OMA): count number of "edits" to change one sequence into another, extensive use in molecular biology
  - Elzinga (2003, 2005): count the number of times the same states occur in the same order in two sequences



## Hamming distance

- Hamming distance is very simple but limited
  - ABABAB and BABABA maximally different temporal rigidity
  - Less of a problem with life course data, typically has long runs in same state
  - Buchmann and Sacchi (1995) factor analyse occupational characteristics to define inter-occupational distances
  - Lesnard (2006) proposes a "dynamic" Hamming distance for time-diary data: distances defined by transition rates between activities, varying throughout the day
  - If time has a meaningful ruler (as with daily time use) Hamming's time-rigidity is a positive advantage
  - For more elastic "developmental" time, perhaps less appropriate





- Degenne's method promising but untried
- Dijkstra-Taris more or less directly superseded by OMA
- Elzinga's method shows a lot of promise, plus has strong intuitive basis



- Given two sequences, *s*<sub>1</sub> and *s*<sub>2</sub>, drawn from an "alphabet" *S*, the Optimal Matching Algorithm generates a distance measure based on "elementary operations":
  - substitution and
  - insertion and deletion
- Substitution replaces an element in *s*<sub>1</sub> with the corresponding element in *s*<sub>2</sub>
- Insertions and deletions equivalently delete an element in  $s_1$  or insert an element in  $s_2$  (or vice versa); because of the equivalence they are known as *indels*



## The Optimal Matching Algorithm-distance

- The distance between *s*<sub>1</sub> and *s*<sub>2</sub> is defined as the least expensive path from one to the other using the elementary operations
- OMA allows a matrix of costs for all pairwise substitutions, and an indel cost to be specified
- The OM algorithm is a dynamic programming technique which finds the "cheapest" path in a time- and memory-efficient manner



- Andrew Abbott has been the main evangelist for OMA in sociology
- A string of articles from Abbott and Forrest (1986), Abbott and Hrycak (1990) *etc.*to a retrospective review (Abbott and Tsay, 2000) with a related debate (Levine, 2000; Wu, 2000; Abbott, 2000)
- In Abbott (1995a), demonstrates OMA to be more general than Dijkstra–Taris



- OM is becoming widely used, particularly since Götz Rohwer incorporated it in TDA
  - Class careers: Halpin and Chan (1998)
  - Women's careers in finance: Blair-Loy (1999)
  - Transition from school to work: Scherer (2001), McVicar and Anyadike-Danes (2002)
  - Methods paper: Brüderl and Scherer (2004)
  - Male careers: Anyadike-Danes and McVicar (2005)
  - Time use: Wilson (2006)
  - Gendered careers: Levy et al. (2006)
  - Housing, employment, marriage, fertility: Pollock (2007)
- More recently, Kohler et al have released a Stata package, SQ, for OMA



## Problems with OM for lifecourse data

- For sequences that are naturally discrete in time, OM works well
- But life course sequences are have highly variable spell lengths, and the discrete representation may not suit OM as well
- For instance, given  $s_1 = ABBD$ ,  $s_2 = ABCD$ ,  $s_3 = ABDD$ , all three sequences will be equidistant
- But sociologically, *s*<sup>1</sup> and *s*<sup>3</sup> are clearly closer
- OM doesn't recognise the continuity; a slight adjustment of spell length is treated as being as expensive as the introduction of another spell



#### Change the time scale?

- One suggestion (Abbott made it in early work) is to change the time scale: use log time, for instance
- However, this exacerbates the discretisation





- I propose instead a modification of the OM algorithm that has an analogous effect: scale the costs of elementary operations according to the length of the affected spells
- There are limits to what we can change in the algorithm without degrading its performance
- In particular, the algorithm has no memory (*e.g.*, "now deleting an element from this spell for the second time")
- However, we can take account of spell length in setting costs



- I begin by outlining the operation of the OM algorithm
- It uses dynamic programming techniques to efficiently determine the cheapest set of edits to transform one sequence into another hence "optimal"
- Operates by calculating the elements of a matrix where each element  $C_{ij} = min(c_{i-1,j-1} + ! i_{i,j}, c_{i,j-1} + ., c_{i-1,j} + .)$ ,
- ! is the substitution matrix, and the first row and column are filled with the cumulative insertion/deletion costs
- A diagonal move represents a substition, right represents deletion, down represents insertion
- Bottom right cell eventually contains the optimal cost



## Working through OM

Cell value: 
$$min(c_{i-1,j-1} + !_{i,j}, c_{i,j-1} + , c_{i-1,j} + )$$
  
=  $min(0 + 2, 2 + 2, 2 + 2) = 2$   
=  $min(2 + 1, 2 + 2, 4 + 2) = 3$   
=  $min(4 + 0, 3 + 2, 6 + 2) = 4$   
=  $min(6 + 1, 4 + 2, 8 + 2) = 6$ 

А В С D С D *s*<sub>2</sub> А А В 

 $s_1$ 

0	2	4	6	8	
2	2	3	4	6	
4	4	4	4	4	
6	4	5	6	6	
8	6	5	7	8	
10	8	6	6	8	



To convert ABCD into CDAAB the following set of operations gives the cheapest path:

Operation	Intermediate state	Cost
	ABCD	= 0
insert C	CABCD	+2 = 2
insert D	CDABCD	+2 = 4
const A = A	CDABCD	+0 = 4
subs $B \rightarrow A$	CDAACD	+1 = 5
subs $C \rightarrow B$	CDAABD	+1 = 6
delete D	CDAAB-	+2 = 8

Where sequence lengths are variable the distance is often scaled inversely with the longer of the pair: 8 units thus become an pairwise distance of 1.6



#### Efficient

- This strategy calculates the optimum path efficiently
- Memory requirements are proportional to  $l_1 \times l_2$
- Time requirement is  $a + b(l_1 \times l_2)$
- Both are small relative to the  $O(N^2)$  required to process every pairwise comparison between N sequences
- Nonetheless, for realistic data sets the procedure completes acceptably fast
- **However**, adaptations of the algorithm risk destroying its efficiency, for instance by requiring exponentially growing memory or processing time



## Variant OMA

- The modified algorithm I propose treats spell length in the following manner
  - single unit spell costed as per OM
  - total cost of a multiple unit spell to be strictly increasing in length
  - cost of each unit to be strictly decreasing with length of spell
- Costing units at *basecost*  $\times \frac{1}{\sqrt{len}}$  achieves this:
  - Cost of 1 unit spell: *basecost* × 1
  - Cost of 2 unit spell: basecost × (0.707 + 0.707) = basecost × 1.414, etc.
- So will other exponents, other functions



- Insertion and deletion costs are modified directly and equally, as they are equivalent: an insertion in  $s_1$  is equivalent to a deletion in  $s_2$
- Substitution costs must also be modified, as a substitution is a deletion followed by an insertion
- The direction of the substitution is unimportant, so we cost it as if it involved a deletion in the longer subsequence
- If  $s_{1i}$  is unique, and  $s_{2j}$  is part of a run of two, the substitution cost is divided by  $\sqrt{2}$
- If  $s_{1i}$  is one of three and  $s_{2i}$  one of two, division is by  $\sqrt{3}$



#### The matrix operations



0	2	4	6	8
2.0	2.0	3.0	4.0	6.0
4.0	4.0	4.0	4.0	4.0
5.4	4.0	4.7	5.4	5.4
6.8	5.4	4.7	6.1	6.8
8.8	7.4	5.4	5.7	7.7



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Sequences		OMA	New variant		
А	В		exp = 0.5	exp = 0.75	
12343	23414	1.00	1.00	1.00	
12343	22213	1.00	0.83	0.78	
12343	22222	1.00	0.45	0.30	
23414	22213	0.80	0.55	0.46	
23414	22222	1.20	0.54	0.36	
22213	22222	0.40	0.18	0.12	



- Implemented in Stata
- C plugin, platform dependent, fast, relatively simple code
- Platform independent Mata implementation also possible, but very slow by comparison



## Simulated data

- Tests are run with simulated data, 4 states, 20 time-units, with
  - a low probability of multiple observations in sequence (82% of cases have max run of  $\leq$  2)
  - a medium probability (79%  $\leq$  5)
  - and a higher probability (79%  $\leq$  11)
- As expected, the higher the presence of runs the lower the mean scores
- But correlations remain high, even excluding matches between identical sequences (distance zero)



## Correlations

		All	Non-identical
Low	Pearson	0.9978	0.9496
	Spearman	0.9984	0.9420
Med	Pearson	0.9931	0.8092
	Spearman	0.9946	0.7985
High	Pearson	0.9821	0.8128
	Spearman	0.9949	0.8132



## Correlation of OMA and OMAv: low level of runs



#### Correlation of OMA and OMAv: medium level of runs



## Correlation of OMA and OMAv: high level of runs



- But what about real data and the "end product"?
- Typically "empirical typologies" generated by cluster analysis
- Using BHPS fertility and labour market histories, I construct 5-year labour market histories for women: 2 years before and 3 years after a birth
- Encoded as monthly status, 4 states:
  - Full-time employed
  - Part-time employed
  - Unemployed
  - Not in labour market



## OM costings

• Substitution matrix:

FTE PTE UE NonE

	0	1	2	3
	1	0	1	2
	2	1	0	1
Ε	3	2	1	0

• *Indel* cost: 2 units



## Correlation of OMA and OMAv: labour market data



#### Results

- Correlation of 0.97 between OMA and OMAv, excluding identical sequences higher than expected
- Spearman correlation similar



- Fit an 8 cluster solution to both data sets
- By inspection gives an acceptable result
- Very high level of agreement, especially for "low entropy" sequences (*i.e.*, near 100% dominated by a single state)
- Less agreement where more is "going on"



# Comparing the eight-cluster solution

OMAv	OMA					Total			
	1	2	3	4	5	6	7	8	
1 (slide 32)	263	27	0	0	0	2	1	0	293
2 (slide 33)	0	39	7	0	2	0	0	0	48
3 (slide 34)	0	0	18	0	0	0	0	0	18
4 (slide 35)	0	0	19	54	1	0	0	0	74
5 (slide 36)	0	0	0	0	33	0	0	0	33
6 (slide 37)	0	1	0	0	0	21	0	0	22
7 (slide 38)	0	0	0	0	0	3	18	0	21
8 (slide 39)	0	0	27	0	0	0	0	139	166
Total	263	67	71	54	36	26	19	139	675
(41,42)									



## 8 cluster solution – OMA (left) and OMAv (right)





## Cluster 1 – OMA (left) and OMAv (right)





## Cluster 2 – OMA (left) and OMAv (right)





## Cluster 3 – OMA (left) and OMAv (right)





## Cluster 4 – OMA (left) and OMAv (right)





## Cluster 5 – OMA (left) and OMAv (right)





## Cluster 6 – OMA (left) and OMAv (right)





## Cluster 7 – OMA (left) and OMAv (right)





## Cluster 8 – OMA (left) and OMAv (right)





- High agreement
- Two major sources of disagreement
  - 27 cases move from OMA cluster 2 to OMAv cluster 1
  - OMA cluster 3 scattered across OMAv clusters 2, 3, 4 and 8
- The 27 OMA cluster 2 cases are arguably better off in cluster 1 (dominated by non-employment) than cluster 2, where early unemployment is matched to early FT-employment



# OMA Cluster 1/2 split





#### OMA Cluster 3 split





#### OMA Cluster 3 details

- Cases moved to cluster 8 are initially FTE with lots of transitions in and out of employment later, and are moved to the predominant FTE cluster
- Those moved to cluster 4 transition from FTE to PTE around the birth, and largely stay there, and merge with the predominant PTE cluster
- Those that remain in cluster 3 are a distinct group: initially FTE, try to remain in the labour market but finally drop out
- to cluster 4: FTE with late shift to PTE, matched to predominant FTE cluster
- Those that move to cluster 2 are also initially FTE but drop out, typically with a short spell of PTE, and are matched with very similar trajectories that do not have the short PTE spell



## Conclusions

- OMAv and OMV generate very but not completely similar results
- To the extent that they differ, it is arguable that it gives a superior clustering of the more complicated trajectories
- Extent of similarity is greater than I had expected
  - Conventional OMA may be adequate where all sequences characterised by long runs
  - Greater mix of long and short runs may be different result is likely to depend on the data to some degree
- As seen, the differences are greatest in the clusters characterised by higher "entropy" sequences these are the ones most sensitive to the distance measure; even naïve matching (*e.g.*, Hamming) will match the simple cases
- "More research is needed"



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